EB The connected circular metric dimension of a graph

¹S. Sheeja and ²K.Rajendran

¹Research Scholar, Department of Mathematics,

Vels Institute of Science, Technology and Advanced Studies, Pallavaram, Chennai-117. email: sheeja1304@gmail.com

²Associate Professor, Department of Mathematics,

Vels Institute of Science, Technology and Advanced Studies, Pallavaram, Chennai-117. email: gkrajendra59@gmail.com

Abstract

Let G = (V, E) be a simple graph and u, v be any two vertices of G. Then the circular distance between u and v denoted by $D^{c}(u, v)$ and is defined by

 $D^{c}(u, v) = \begin{cases} D(u, v) + d(u, v) & \text{if } u \neq v \\ 0 & \text{if } u = v \end{cases}$

where D(u, v) and d(u, v) are detour distance and distance between u and v respectively. Let $W = \{w_1, w_2, ..., w_k\} \subset V(G)$ and $v \in V(G)$. The representation cr(v/W) of v with respect to W is the k-tuple $(D^c(v, w_1), D^c(v, w_2), ..., D^c(v, w_k))$. Then W is called a circular resolving set if different vertices of G have different representations with respect to W. A circular resolving set W is called connected circular resolving set, G[W] is connected. The minimum cardinality of a connected circular resolving set in a graph G is its connected circular metric dimension of G and is denoted by $cdim_c(G)$. The connected circular metric dimension of some standard graphs are determined. Some general properties satisfies by this concept are studied. Connected graphs of order $n \ge 3$ with connected circular metric dimension 1 are characterized. Necessary condition for the connected circular metric dimension to be n - 1 is given.

Keywords: distance, detour distance, circular distance, circular resolving set, connected circular metric dimension.

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Section A-Research paper

1. Introduction and Preliminaries

Let G be a simple graph with vertex set V(G) and edge set E(G). The order of a graph G is |V(G)|, its number of vertices denoted by n. The size of a graph G is |E(G)|, its number of edges denoted by m. For basic graph theory terminology, we refer [3]. The degree, deg(v) of a vertex $v \in V(G)$ is the number of edges incident to v. We denote $\Delta(G)$ the maximum degree of a graph G. The distance d(u, v) between two vertices $u, v \in V(G)$ is the length of a shortest path between them. These concepts were studied in [1,2,5,10-14,17-24,26,27]. The detour *distance* D(u, v) between two vertices $u, v \in V(G)$ is the length of a longest path between them. These concepts were studied in [6-8,15]. Let $W = \{w_1, w_2, \dots, w_k\} \subset V(G)$ and $v \in V(G)$. The representation r(v/W) of v with respect to W is the k-tuple $(d(v, w_1), d(v, w_2), \dots, d(v, w_k))$. Then W is called a resolving set if different vertices of G have different representations with respect to W. A resolving set of minimum number of elements is called a *basis* for G and the cardinality of the basis is known as the *metric* dimension of G, represented by dim(G). These concepts were studied in [4,28]. The Dr(v/W)of W representation v with respect to is the *k*-tuple $(D(v, w_1), D(v, w_2), \dots, D(v, w_k))$. Then W is called a *detour resolving set* if different vertices of G have different representations with respect to W. A detour resolving set of minimum number of elements is called a *detour basis* for G and the cardinality of the basis is known as the *detour metric dimension* of *G*, represented by *Ddim*(*G*).

The circular distance between u and v is denoted by $D^{c}(u, v)$ and is defined by

$$D^{c}(u,v) = \begin{cases} D(u,v) + d(u,v) & \text{if } u \neq v \\ 0 & \text{if } u = v \end{cases}$$

An u - v path of length $D^{c}(u, v)$ is called a u - v circular. The circular diameter is the maximum circular distance be a pair of vertices of G. It is denoted by the $D^{c}(G)$. A circular path of length $D^{c}(G)$ is called the circular diametral path. These concepts were studied in [16].

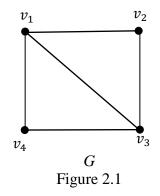
Let $W = \{w_1, w_2, ..., w_k\} \subset V(G)$ and $v \in V(G)$. The representation cr(v/W) of vwith respect to W is the k-tuple $(D^c(v, w_1), D^c(v, w_2), ..., D^c(v, w_k))$. Then W is called a *circular resolving set* if different vertices of G have different representations with respect to W. A circular resolving set of minimum number of elements is called a circular basis for G and the cardinality of the basis is known as the *circular metric dimension* of G, represented by cdim(G). These concepts were studied in [25]. In this article, we study a new metric dimension called the connected circular metric dimension of a graph. The following theorem is used in the sequel.

Theorem 1.3. [28]] For the complete graph $G = K_n$ $(n \ge 2)$, cdim(G) = n - 1.

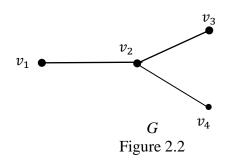
2. The connected circular metric dimension of a graph

Definition 2.1. A circular resolving set W is called a *connected circular resolving set* of G if G[W] is connected. The minimum cardinality of a connected circular resolving set in a graph G is its *connected circular metric dimension* of G and is denoted by $cdim_c(G)$.

Example 2.2. For the graph G given in Figure 2.1, let $W = \{v_1, v_2\}$. Then $cr(v_1/W) = (0, 4)$, $cr(v_2/W) = (4,0)$, $cr(v_3/W) = (3,4)$, $cr(v_4/W) = (4,5)$. Since cr(v/W) are distinct for all $v \in V(C_n)$, it follows that W is a circular resolving set of G. Since G[W] is connected, W is a connected circular resolving set of G and so $cdim_c(G) \leq 2$. Also since no singleton subset of V(G) is a circular resolving set of G, we have $cdim_c(G) = 2$.



Example 2.3. For the graph *G* given in Figure 2.2, no singleton subset of V(G) is a circular resolving set of *G*, we have $cdim_c(G) \ge 2$. Let $W = \{v_1, v_3\}$. Then $cr(v_1/W) = (0,4)$, $cr(v_2/W) = (2,2)$, $cr(v_3/W) = (4,0)$, $cr(v_4/W) = (4,4)$. Since cr(v/W) are distinct for all $v \in V(C_n)$, it follows that *W* is a circular resolving set of *G*. Since *G*[*W*] is not connected, *W* is not a connected circular resolving set of *G*. It is easily verified that no two-element subset of V(G) is not a circular resolving set of *G* and so $cdim_c(G) \ge 3$. Let $W_1 = \{v_1, v_2, v_3\}$. Then W_1 is a connected circular resolving set of *G* so that $cdim_c(G) = 3$.



Observation 2.4. (i) Let *G* be a connected graph of order $n \ge 2$. Then $1 \le cdim_c(G) \le n - 1$. (ii) Each cut vertex of *G* belongs to every connected resolving set of *G*.

In the following we determine the connected circular metric dimension of some standard graphs.

Theorem 2.5. For the graph $G = P_n$ $(n \ge 2)$, $cdim_c(G) = 1$.

Proof. Let $V(P_n) = \{v_1, v_2, ..., v_n\}$ and let $W = \{v_1\}$. Then $D^c(v_1, v_i) = 2(i - 1), (1 \le i \le n)$. Since cr(v/W) is distinct for all $v \in V(P_n)$, it follows that W is a circular resolving set of G. Also, G[W] is connected, Hence W is a connected circular resolving set of G so that $cdim_c(G) = 1$.

Theorem 2.6. For the cycle $G = C_n$, $n \ge 3$, $cdim_c(G) = n - 1$.

Proof. Let $V(G) = \{v_1, v_2, ..., v_n\}$ and let $W = \{v_1, v_2, ..., v_{n-1}\}$, The circular representations of (n-1) tuples are as follows

$$cr(v_1/W) = (0, n, n, ..., n),$$

$$cr(v_2/W) = (n, 0, n, n, ..., n),$$

$$cr(v_3/W) = (n, n, 0, n, n, ..., n),$$

 $cr(v_{n-1}/W) = (n, n, n, ..., n, 0)$ $cr(v_n/W) = (n, n, n, n, ..., n).$

Since cr(v/W) are distinct for all $v \in V(C_n)$, it follows that W is a circular resolving set of G. Since G[W] is connected, W is a connected circular resolving set of G. Therefore $cdim_c(G) \le n-1$. We substantiate that $cdim_c(G) = n - 1$.Consider, however, that $cdim_c(G) \le n-2$. 1735

Eur. Chem. Bull. 2023, 12(Special Issue 7), 1732-1744

Then, a set S'exists such that $|S'| \le n-2$. As a result, there are at least two vertices u, v that satisfy the contradiction cr(u/S') = cr(v/S') = (n, n, ..., n). Consequently, $cdim_c(G) = n-1$.

Theorem 2.7. For the complete graph $G = K_n$, $n \ge 2$, $cdim_c(G) = n - 1$.

Proof. The proof is similar to the Theorem 2.6.

Theorem 2.8. For the wheel graph $G = W_n$, $n \ge 3$, $cdim_c(G) = n - 4$.

Proof. Let $V(G) = \{v_1, v_2, ..., v_{n-1}, u\}$ and $W = \{v_1, v_2, ..., v_{n-4}\}$, Then the circular representations of (n - 4) tuples are as follows

$$cr(v_1/W) = (0, n, n + 1, ..., n + 1, n + 1)$$

$$cr(v_2/W) = (n, 0, n, n + 1, ..., n + 1, n + 1)$$

$$cr(v_3/W) = (n + 1, n, 0, n, n + 1, ..., n + 1)$$

$$\begin{split} cr(v_{n-4}/W) &= (n+1,n+1,n+1,\dots,n,0)\\ cr(v_{n-3}/W) &= (n+1,n+1,n+1,\dots,n+1,n)\\ cr(v_{n-2}/W) &= (n+1,n+1,n+1,\dots,n+1,n+1)\\ cr(v_{n-1}/W) &= (n,n+1,n+1,\dots,n+1,n+1)\\ cr(u/W) &= (n,n,n,n,\dots,n,n). \end{split}$$

Since cr(v/W) are distinct for all $v \in V(W_n)$, it follows that W is a circular resolving set of G. Since G[W] is connected, W is a connected circular resolving set of G. Therefore $cdim_c(G) \le n - 4$. We substantiate that $cdim_c(G) = n - 4$. Consider, however, that $cdim_c(G) \le n - 5$. Then, a set W'exists such that $|W'| \le n - 5$. As a result, there are at least two vertices u, v that satisfy the contradiction

$$cr(u/S') = cr(v/S') = (n + 1, n + 1, ..., n + 1).$$

Consequently, $cdim_c(G) = n - 4$.

Theorem 2.9. For the complete bipartite graph $G = K_{r,s}$, $(1 \le r \le s)$,

$$cdim_{c}(G) = \begin{cases} 1; r = 1, 1 \le s \le 2, \\ r + s - 2; r = 1, s \ge 3. \\ r + s - 1; 2 \le r \le s \end{cases}$$

Proof. Let $X = \{x_1, x_2, ..., x_r\}$ and $Y = \{y_1, y_2, ..., y_s\}$ be the two bipartite sets of *G*. We have the three cases.

Case (i): $r = 1, 1 \le s \le 2$. The result follows from Theorem 2.5.

Case (ii): $r = 1, s \ge 3$.

Let $W = V(G) - \{x_1, y_s\}$. Then the circular metric representations (n - 2) tuples are as follows:

$$cr(x_1/W) = (2,2,2,...,2,2)$$

 $cr(y_1/W) = (0,4,4,...,4,4)$
 $cr(y_2/W) = (4,0,4,...,4,4)$

$$cr(y_{s-1}/W) = (4,4,4,\dots,4,0)$$

 $cr(y_s/W) = (4,4,4,\dots,4,4).$

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Since the representation are distinct and G[W] is connected, W is a connected circular resolving set of G so that $cdim_c(G) \le r + s - 2$. We demonstrate that $cdim_c(G) = r + s - 2$. On the other hand, imagine that $cdim_c(G) \le r + s - 3$. Then there exists a circular resolving set W' such that. $|W'| \le r + s - 3$. As a result, there are at least two end vertices $u, v \in V \setminus W'$ such that cr(u/W') = cr(v/W') = (4,4,4,...,4,4), which is incoherent. As a result, $cdim_c(G) = r + s - 2$.

Cases (iii): $2 \le r \le s$.

Let $W = V(G) - \{y_s\}$. Then the circular metric representations (r + s - 1) tuples are as follows:

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$$cr(x_1/W) = (0, r + s - 1, r + s - 1, ..., r + s - 1)$$

 $cr(x_2/W) = (r + s - 1, 0, r + s - 1, ..., r + s - 1)$

$$cr(x_r/W) = (r + s - 1, r + s - 1, ..., 0, r + s - 1, ..., r + s - 1)$$

$$r^{th} place$$

$$cr(y_1/W) = (r + s - 1, r + s - 1, ..., r + s - 1, 0, r + s - 1 ..., r + s - 1)$$

$$(r + 1)^{th} place$$

$$cr(y_2/W) = (r + s - 1, r + s - 1, ..., r + s - 1, 0, r + s - 1 ..., r + s - 1)$$

$$(r + 2)^{th} place$$

$$cr(y_{s-1}/W) = (r + s - 1, r + s - 1, r + s - 1, ..., r + s - 1, r + s - 1, ..., 0)$$

Eur. Chem. Bull. 2023, 12(Special Issue 7), 1732-1744

Section A-Research paper

$$(r + s - 1)^{\text{th}}$$
 place
 $cr(y_s/W) = (r + s - 1, r + s - 1, r + s - 1, ..., r + s - 1, r + s - 1, ..., r + s - 1)$

Since the representation are distinct and G[W] is connected, W is a connected circular resolving set of G so that $cdim_c(G) \le r + s - 1$. We demonstrate that $cdim_c(G) = r + s - 1$. Consider however, that $cdim_c(G) \le r + s - 2$. If so, a circular resolving set W' exists such that $|W'| \le r + s - 2$ As a result, there are at least two vertices, $u, v \in V \setminus W'$ such that cr(u/W') = cr(v/W') = (r + s - 1, r + s - 1, r + s - 1, ..., r + s - 1), which is incoherent. As a result, $cdim_c(G) = r + s - 1$.

Theorem 2.10. Let *G* be the graph obtained from $K_{1,n-1}$, $(n \ge 3)$, by subdividing the end edges exactly once. Then $cdim_c(G) = n - 1$.

Proof. Let *x* be the central vertex of $K_{1,n-1}$ $(n \ge 4)$ and $\{v_1, v_2, ..., v_{n-1}\}$ be the set of end vertices of *G*. *G* is the graph obtained form $K_{1,n-1}$, $(n \ge 4)$, by subdividing xv_i $(1 \le i \le n - 1)$ by u_i $(1 \le i \le n - 1)$. Let $W = \{x, u_1, u_2, ..., u_{n-2}\}$. Then

$$cr(x/W) = (0,2,2,...,2,2)$$

 $cr(u_1/W) = (2,0,4,4,...,4,4)$
 $cr(u_2/W) = (2,4,0,4,...,4,4)$

$$cr(u_{n-3}/W) = (2,4,4,4,\dots,0,4,4)$$

$$cr(u_{n-2}/W) = (2,4,4,4,\dots,4,4,0)$$

$$cr(u_{n-1}/W) = (2,4,4,\dots,4,4,4)$$

$$cr(v_1/W) = (4,2,6,6,6,\dots,6,6)$$

$$cr(v_2/W) = (4,6,2,6,6,\dots,6,6)$$

$$\begin{split} cr(v_{n-3}/W) &= (4,6,6,\ldots,2,6)\\ cr(v_{n-2}/W) &= (4,6,6,6,\ldots,6,2)\\ cr(v_{n-1}/W) &= (4,6,6,6,\ldots,6,6,). \end{split}$$

Due to the distinctness of the representations, W is a circular resolving set of G. Also G[W] is connected, W is a connected circular resolving set of G so that $cdim_c(G) \le n-1$. We

substantiate that $cdim_c(G) = n - 1$. Consider, however, that $cdim_c(G) \le n - 2$. If so, a circular resolving set W' exists such that $|W'| \le n - 2$ and G[W'] is disconnected. Consequently, $cdim_c(G) = n - 1$.

Theorem 2.11. Let *G* be the graph obtained from C_n , $(n \ge 3)$, by subdividing the edges exactly once. Then $cdim_c(G) = 2n - 1$.

Proof: Let $V(C_n) = \{v_1, v_2, ..., v_n\}$ and $\{u_1, u_2, ..., u_n\}$ be the subdivided vertices of C_n . Then *G* is a cycle contains 2n vertices. By Theorem 2.6, $cdim_c(G) = 2n - 1$.

Theorem2.12. For the crown graph $G = H_{n,n}$, $n \ge 3$, $cdim_c(G) = n$.

Proof. Let $V(G) = \{u_1, u_2, \dots, u_n, v_1, v_2, \dots, v_n\}$ and $E(G) = \{(u_i, v_j); 1 \le i, j \le n; i \ne j\}$,

Let $W = \{u_1, u_3, u_4, \dots, u_n, v_2\}$. Then the circular representations of n tuples are as follows

$$cr(u_1/W) = (0, n, n, n, ..., n, n, n, n)$$

$$cr(u_2/W) = (n, n, n, n, ..., n, n, n, n + 2)$$

$$cr(u_3/W) = (n, 0, n, n, ..., n, n, n, n)$$

$$cr(u_4/W) = (n, n, 0, n, ..., n, n, n, n)$$

$$cr(u_{n-1}/W) = (n, n, n, n, ..., n, 0, n, n)$$

$$cr(u_n/W) = (n, n, n, n, ..., n, n, 0, n)$$

$$cr(v_1/W) = (n + 2, n, n, n, ..., n, n, n, n)$$

$$cr(v_2/W) = (n, n, n, n, ..., n, n, n, n)$$

$$cr(v_3/W) = (n, n + 2, n, n, ..., n, n, n, n)$$

$$cr(v_4/W) = (n, n, n + 2, n, ..., n, n, n, n)$$

$$cr(v_{n-1}/W) = (n, n, n, n, ..., n, n + 2, n, n)$$

 $cr(v_n/W) = (n, n, n, n, ..., n, n, n + 2, n)$

Since cr(v/W) are distinct for all $v \in V(H_{n,n})$, it follows that W is a circular resolving set of G. Since G[W] is connected, W is a connected circular resolving set of G. Therefore $cdim_c(G) \le n$. We substantiate that $cdim_c(G) = n$. Consider, however, that $cdim_c(G) \le n - 1$. Then, a set W'exists such that $|W'| \le n - 1$. As a result, there are at least two vertices u, v that satisfy the contradiction

$$cr(u/S') = cr(v/S') = (n, n, n, n, ..., n, n).$$

Consequently, $cdim_c(G) = n$.

3.Some results on connected circular metric dimension of a graph

Theorem 3.1. For connected graph of order $n \ge 2$, $1 \le cdim(G) \le cdim_c(G) \le n - 1$.

Proof: Any circular resolving set of *G* needs atleast one vertex and so $cdim(G) \ge 1$. Since any connected circular resolving set is also a circular resolving set of *G*, we have $cdim(G) \le cdim_c(G)$. Also since V(G) - x is a connected resolving set of *G*, where $x \in V(G)$ is not a cut vertex of *G*, we have $cdim_c(G) \le n - 1$. Thus $1 \le cdim(G) \le cdim_c(G) \le n - 1$.

Remark 3.2. The bounding bound in Theorem 3.1 bounds are sharp.

For $G = P_n$, $n \ge 2$, by theorem 2.5 $cdim_c(G) = 1$.

For the cycle $G = C_4$, by theorem 2.6 $cdim_c(G) = 3$ and for $G = K_n$, $n \ge 3$, $cdim_c(G) = n - 1$.

Remark 3.3. Also, the bounds in Theorem 3.1 can be strict. For the star $G = K_{1,4}$, cdim(G) = 3 $cdim_c(G) = 4$ and n = 5. Thus $1 < cdim(G) < cdim_c(G) < n - 1$.

Theorem 3.4. Let *G* be a connected graph of order $n \ge 3$ has connected circular metric dimension 1 if and only if $G = P_n$.

Proof. Let $G = P_n$. Then the result follows from Theorem 2.5. Conversely, assume that $cdim_c(G) = 1$. Let $W = \{v\}$ be a minimum connected circular resolving set of G. Then $cr(u/W) = D^c(u, v)$ is a non-negative integer less than 2(n - 1) for each $u \in V(G)$. There exists a vertex $u \in V(G)$ such that d(u, v) = n - 1. This is because the representation of V(G) with regard to W are distinct. As a result, the circular diameter of G is 2(n - 1), implies that $G = P_n$.

Theorem 3.5. Let *G* be a connected graph of order $n \ge 3$. If every pair of vertices of *G* is a circular diametral path of *G*. Then $cdim_c(G) = n - 1$

Proof: Assume that every pair of vertices of *G* is circular diametral path of *G*. Therefore $D^{c}(u, v) = n$ for all $u, v \in V(G)$. Hence it follows that every circular resolving set of *G* contains at least n-1 elements. Also, G[W] is connected. Hence $cdim_{c}(G) = n - 1$.

Remark 3.6. The converse of the Theorem 3.5. need not be true. For the graph $G = K_{1,n-1}$, $cdim_c(G) = n - 1$. But there are at least two vertices say x and y in G such that x - y is not a circular diametral path of G.

Theorem 3.7. For any pair of integers *a* and *b* with $2 \le a \le b$, there exists a connected graph *G* such that cdim(G) = a and $cdim_c(G) = b$.

Proof: For a = b, let $G = K_{a+1}$ then by Theorems 1.1 and 2.7, $cdim(G) = cdim_c(G) = a$. So let $2 \le a < b$. Let P_{b-a} be a path of order b-a+1 and let $V(P_{b-a+1}) = \{v_1, v_2, \dots, v_{b-a+1}\}$. Let G be the graph obtained from by adding the new vertices u_1, u_2, \dots, u_a and introducing the edge $v_{b-a+1}u_i$ $(1 \le i \le a)$. The graph is shown in Figure 3.1.

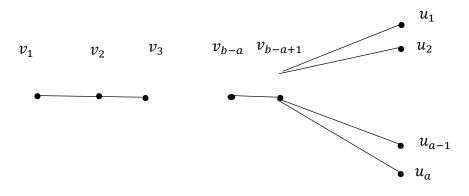


Figure 3.1

First, we prove that cdim(G) = a. Let $Z = \{u_1, u_2, ..., u_a\}$. Then every circular resolving set of *G* contains at least a - 1 vertices from *Z* and the vertex v_1 and so $cdim(G) \ge a - 1 + 1 = a$. Let $S = Z \cup \{v_1\}$. Then *S* is a circular resolving set of *G* so that cdim(G) = a. Next, we prove that $cdim_c(G) = b$. By Observation 2.4(ii), $Z_1 = \{v_2, v_3, ..., v_{b-a+1}\}$ is a subset of every connected circular resolving set of *G*. Also it is easily seen that every connected circular resolving set of *G* contains at least a-1 vertices from *Z* and the vertex v_1 and so $cdim_c(G) = b - a + 1 + a - 1 = b$. Let $S_1 = S \cup Z_1$. Then S_1 is a connected circular resolving set of *G* so that $cdim_c(G) = b$.

Conclusion

This article established a novel circular distance metric called the connected circular metric dimension in graphs. We will develop this concept to incorporate more distance considerations in a subsequent investigation.

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Section A-Research paper

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